

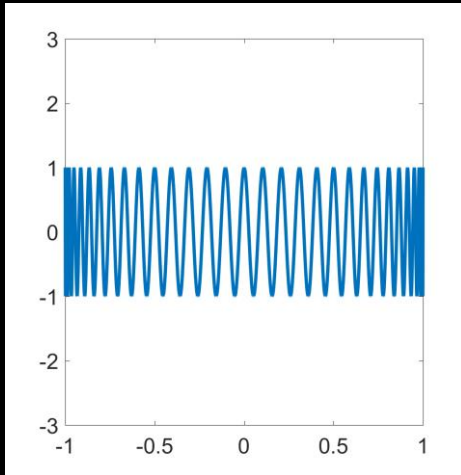
Three Perspectives on Orthogonal Polynomials

Paul Valiant

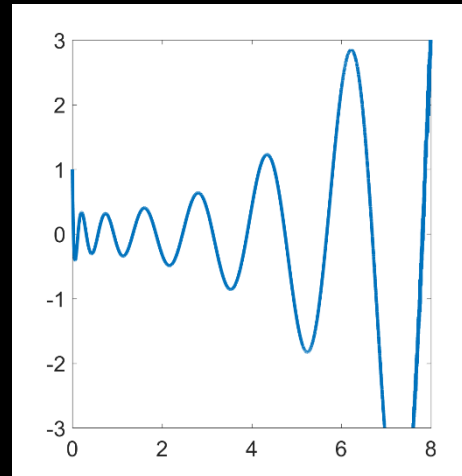
Brown University

(Based on joint work with Gregory Valiant, mostly STOC'11:
“Estimating the Unseen: An $n/\log(n)$ -sample Estimator for Entropy and Support Size, Shown Optimal via New CLTs”)

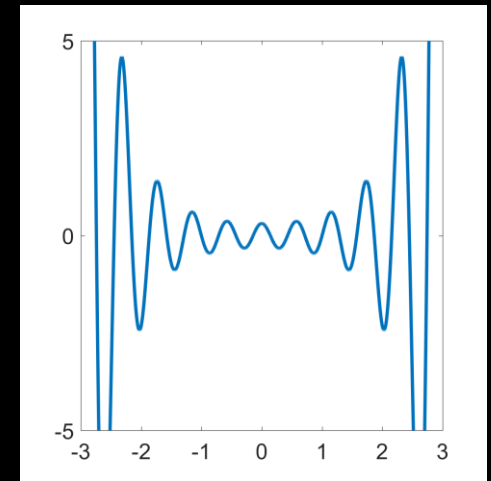
Structure of this talk: 3 polynomial challenges... and solutions



Chebyshev



Laguerre



Hermite

Seems like:
“cosine”

“cosine times
exponential”

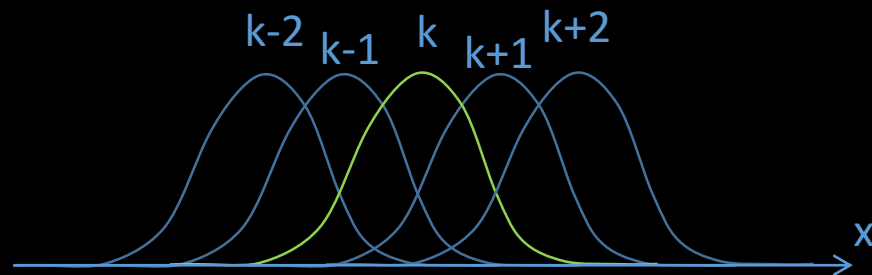
“cosine over
Gaussian”

“Seems,” madam? Nay, *it is*. I know not “seems.” — Hamlet

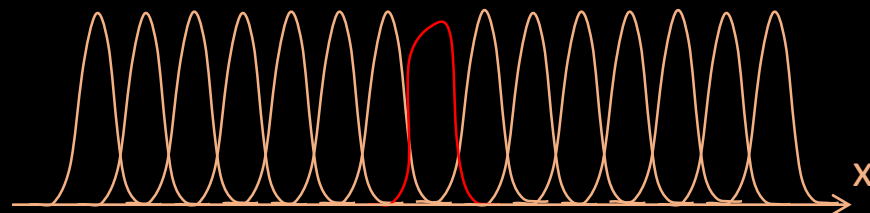




Challenge 1: Poisson bumps \rightarrow thinnest bumps

$$poi(\lambda, k) = \frac{\lambda^k e^{-\lambda}}{k!}$$



Linear transform \downarrow ?



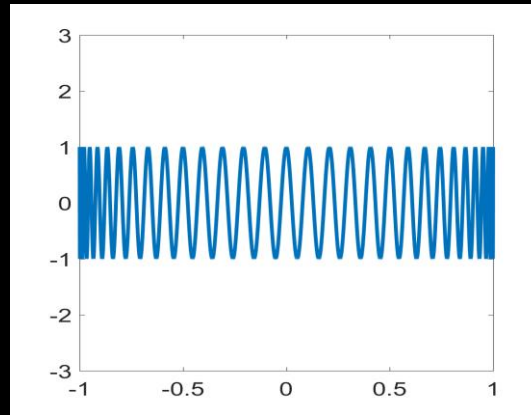
Motivation: Given an event with probability p , $poi(pn, k)$ captures the probability of it occurring exactly k times in $Poi(n)$ samples. Let F_k be total number of events that were observed k times. F_k captures probabilities from . Is there a linear combination of F_k that captures ?

1. Thin ^{x-resolution} as possible
(with bounded coeffs)

2. $\sum \text{bump} \approx 1$
^{y-resolution}

Thm: general log n factor improvement in resolution, #samples

Chebyshev Polynomials



$$T_j(\cos x) = \cos(jx)$$

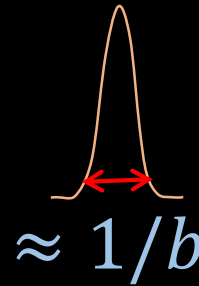
Chebyshev is exactly like cosine,
except on distorted x-axis

- Both unchanged under
x-axis distortion!
- 1. Thin as possible ^{x-resolution}
(with bounded ~~coeffs~~)
 - 2. $\sum_j \text{bump} = 1$ ^{y-resolution}

New question: thin cosine bumps

Thinnest Cosine Bumps

Thinnest linear combination of $\cos(jx)$ for $j < b$:



(Intuition: Fourier transform of degree b gives resolution $1/b$)

1. Thin as possible

(with bounded ~~coeffs~~)

$$2. \sum_j = 1$$

y-resolution

Sum of all possible x-translated bumps is constant



(Trig functions are well-behaved under x-translation)

Chebyshev Takeaways:

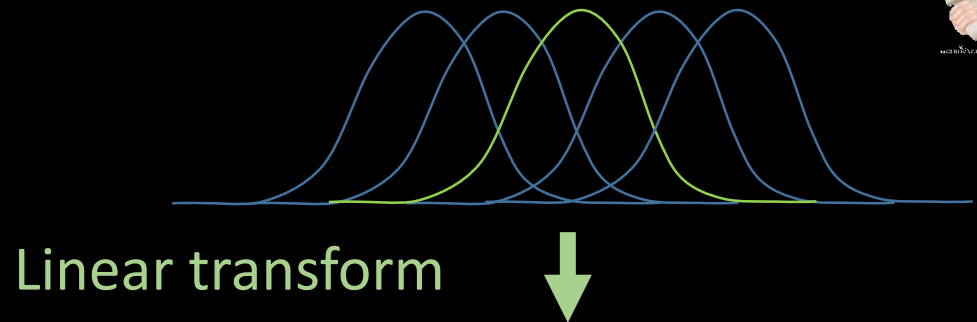
(Modulo x-axis distortion) “polynomials are cosines”



$$poi(\lambda, k) = \frac{\lambda^k e^{-k}}{k!}$$

Motivation: Given an event with probability p , $poi(pn, k)$ captures the probability of it occurring exactly k times in $Poi(n)$ samples. Let F_k be total number of events that were observed k times. F_k captures probabilities from . Is there a linear combination of F_k that captures .

Thm: general $\log n$ factor improvement in resolution, #samples



1. Thin as possible b times thinner

(with coeffs $\ll n$)

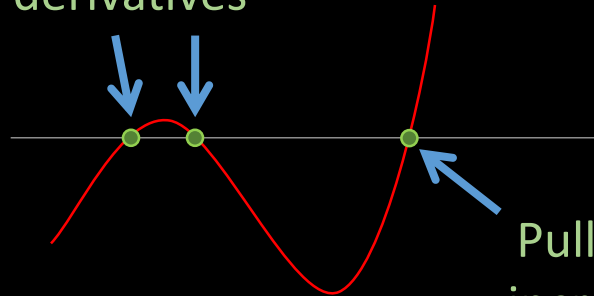
2. $\sum \text{curve} \approx 1$ exp(b)

Thus: $b = \theta(\log n)$

Challenge 2: Exponentially Growing Derivatives

Find: Degree j polynomial with roots at $\varepsilon, 2\varepsilon$; and all remaining roots have much larger derivative, growing exponentially with x

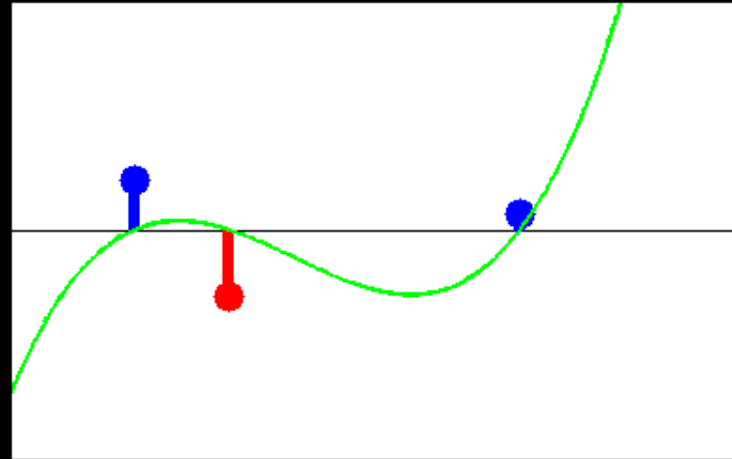
Roots close together have small derivatives



Pulling a root farther away increases its derivative, but only polynomially

Success requires a delicate balancing act!

Orthogonal to Polynomials



Motivation: Want to construct a pair of distributions g^+, g^- that are, respectively, close to the uniform distributions on T and $2T$ elements, but where for each (small) k , the expected number of domain elements seen k times from $Poi(n)$ samples is identical for g^+, g^- . Essentially: find a signed measure $g(x)$ that is

- 1) Orthogonal to $poi(x, k) = \frac{x^k e^{-k}}{k!}$ for each small k ,
- 2) Has most of its positive mass at $1/T$ and most of its negative mass at $1/(2T)$

Fact: If P is a degree j polynomial with distinct real roots $\{x_i\}$, then the signed measure h_p having point mass $1/P'(x_i)$ at each root x_i is orthogonal to all polynomials of degree $\leq j-2$

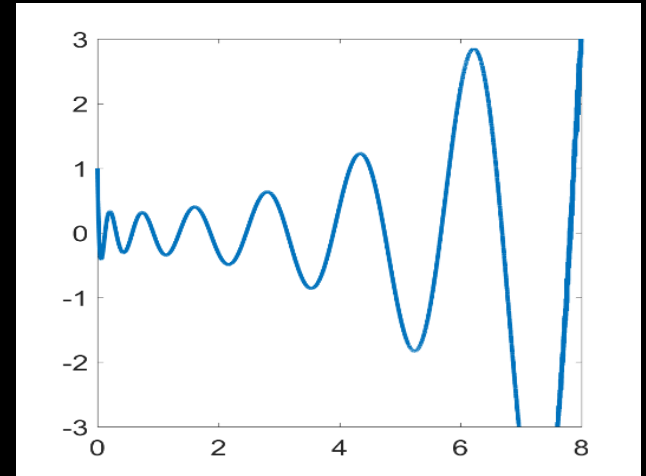
$$g(x) \triangleq e^x h(x)$$

- Essentially: find a signed measure $h(x)$ that is
- 1) Orthogonal to all degree $\leq k$ polynomials
 - 2) Has most of its positive mass at $1/T$ and most of its negative mass at $1/(2T)$
 - and otherwise decays $\ll e^{-x}$

Task: find P such that $P'(x_i)$ grows exponentially in x_i

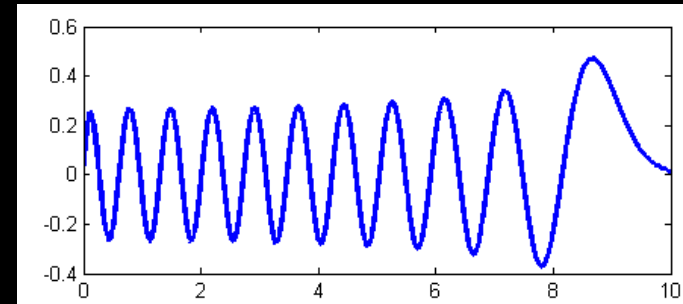
Laguerre Polynomials

Defined by $L_n(x) = e^x \frac{d^n}{dx^n} \frac{e^{-x} x^n}{n!}$ and
 orthogonal as: $\int_0^\infty L_n(x) L_m(x) e^{-x} dx = [m = n]$



Why should the derivative be so nicely behaved at its roots, in particular, growing exponentially?

Transform the Laguerre: $v = e^{-x^2/2} \sqrt{x} \cdot L_n(x^2)$



Many differential equations, including

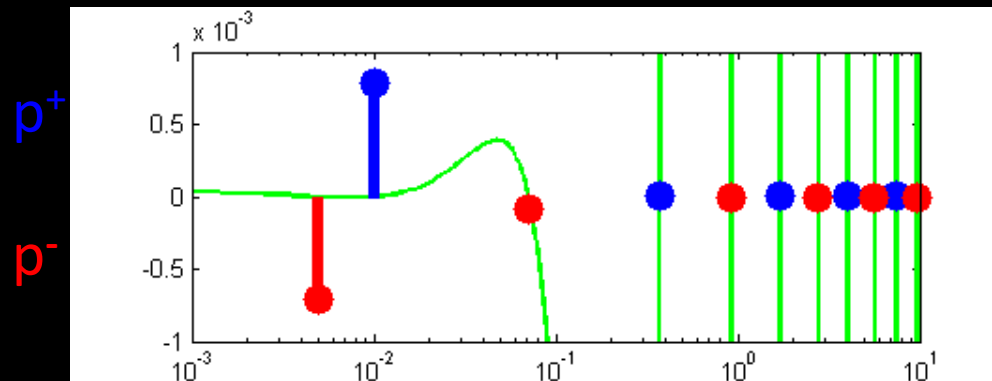
$$v'' + \left(4n + 2 - x^2 + \frac{1}{4x^2}\right) v = 0 \quad \text{Almost harmonic motion, } v \rightarrow \text{sine}$$

Nicely spaced zeros, and max derivative at the zeros

The Construction

Recall: We want a *signed measure* g on the positive reals that:

- Is orthogonal to low degree polynomials
- Decays exponentially fast
- Its positive portion has most of its mass at 2ϵ
- Its negative portion has most of its mass at ϵ



Theorem: p^+ is “close” to $U_{n/2}$, and p^- is “close” to U_n , and p^+ and p^- are indistinguishable via $cn/\log n$ samples

(Modulo diff-eq distortion) “polynomials are $e^x \sin(x)$ ”



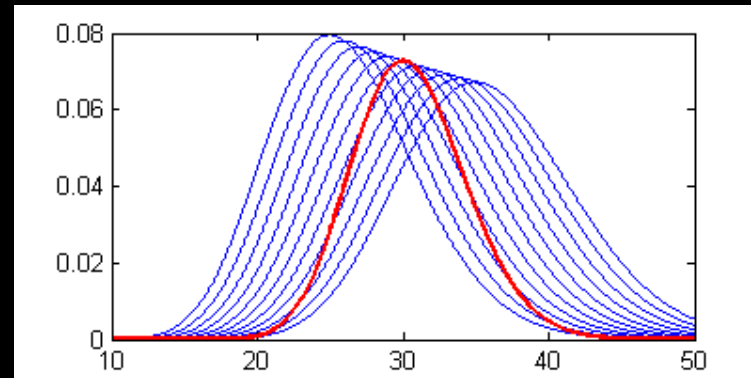
Challenge 3: exponentially good bump approximations

Motivation: Previously, constructed lowerbound distributions g^+, g^- where expectation of every measurement matched. Lower bound? No... until we show variances match too.

Aim: show that variances can be approximated as linear combinations of expectations, with moderate coefficients; thus matching means implies matching variances. Since means come from $\text{poi}(j, x)$, second moments come from $\text{poi}(j, x)^2$.

Find a linear combination over j of $\text{poi}(x, j)$ that approximates $\text{poi}(x, k)^2$ to within ε , using coefficients $\leq 1/\varepsilon$

Think of $\varepsilon = 1/\exp(j)$



These look like Gaussians!

- 1) What's the answer for Gaussians?
- 2) Analyze via Hermite polynomials instead

Approximating “Thin” Gaussians as Linear Combinations of Gaussians

What do we convolve a Gaussian with to approximate a *thinner* Gaussian?

(Other direction is easy, since convolving Gaussians adds their variances)

“Blurring is easy, unblurring is hard” → can only do it approximately

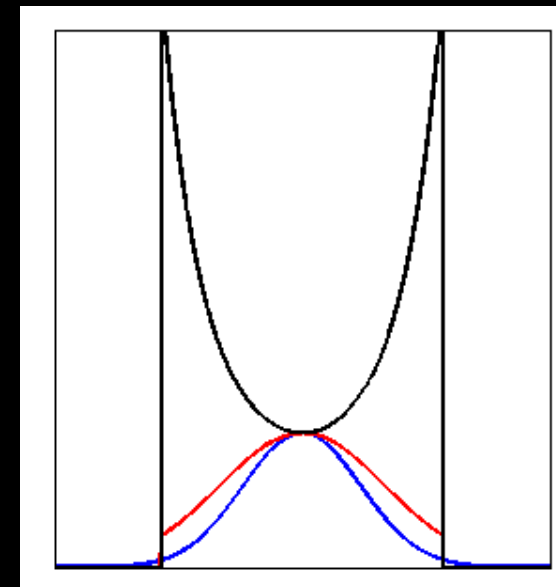
How to analyze? **Fourier transform!** Convolution becomes multiplication

Now: what do we *multiply* a Gaussian with to approximate a *fatter* Gaussian?

$$e^{-x^2} \cdot ??? = e^{-x^2/2}$$

$$e^{x^2/2} \quad \text{Problem: blows up}$$

Answer: if we want to approximate to within ϵ , we only need to approximate out to where $e^{-x^2/2} = \epsilon$. How big is $e^{x^2/2}$ here? $1/\epsilon$



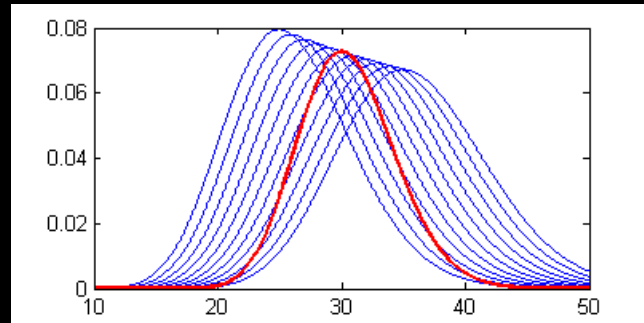
Result: Can approximate to within ϵ using coefficients no bigger than $1/\epsilon$

Hermite Polynomials



Poissons seem a lot like Gaussians

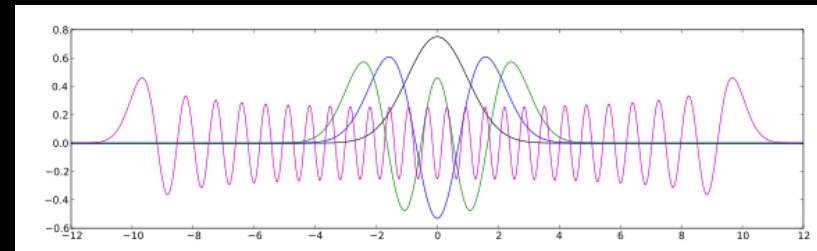
(I was stuck here for about a month)



Idea: $x \rightarrow y^2$

$$\frac{x^j e^{-x}}{j!} = \frac{y^{2j} e^{-y^2}}{j!} \xrightarrow{\text{fourier}} \frac{1}{j!} \frac{d^{2j}}{dw^{2j}} e^{-w^2} = \frac{1}{j!} H_{2j}(w) e^{-w^2}$$

Hermite Polynomials! Sequence of orthogonal polynomials, from which we take the even ones: orthogonal basis for even functions

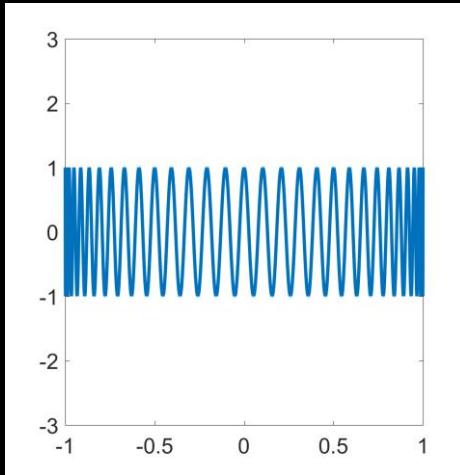


To express **any function** in this basis, just compute each coefficient as an **inner product**

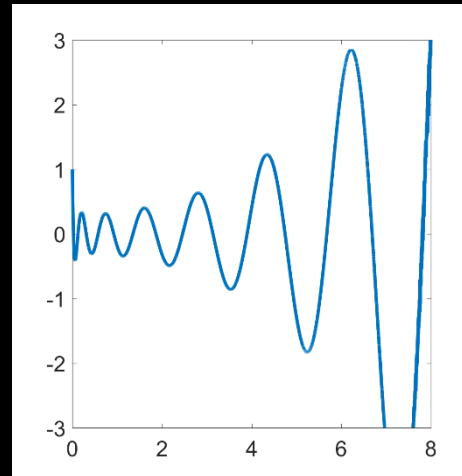
Which function? Fourier transform of “thin” Poisson, cut off at ϵ

Proposition: Can approximate $\Pr[\text{Poi}(2\lambda) = k]$ to within ϵ as a linear combination $\sum_j \alpha_{k,j} \Pr[\text{Poi}(\lambda) = j]$ with coefficients that sum to $\sum_j |\alpha_{k,j}| \leq \frac{1}{\epsilon} 200 \max\{\sqrt[4]{k}, 24 \log^{\frac{3}{2}} \frac{1}{\epsilon}\}$

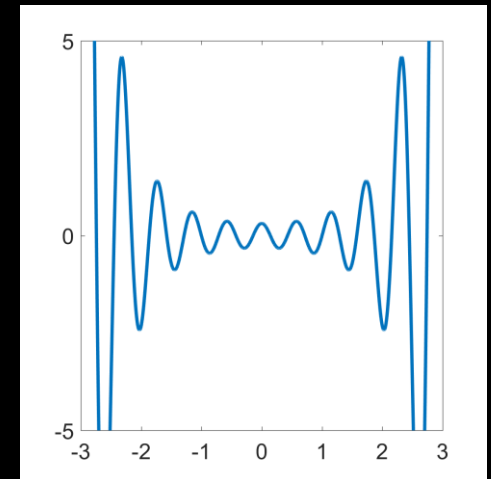
Structure of this talk: 3 polynomial challenges... and solutions



Chebyshev



Laguerre



Hermite

Seems like:

“cosine”

“cosine times
exponential”

“cosine over
Gaussian”

“Seems,” madam? Nay, *it is*. I know not “seems.” — Hamlet

